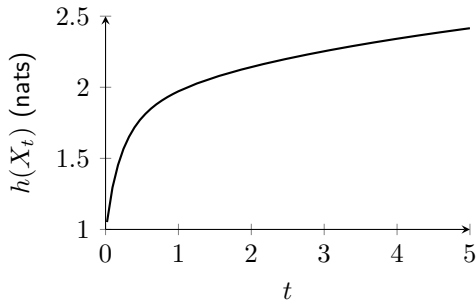
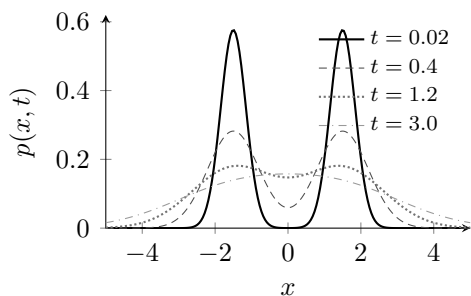


Entropy in Discrete and Continuous Spaces and a Relation to the Heat Equation

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Euler Circle IRPW, July 9, 2026

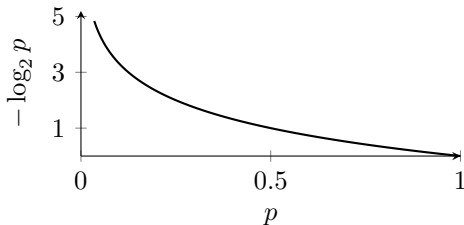


Surprise

k digits name 2^k outcomes N equally likely $\Rightarrow \log_2 N$ digits

a bit is one binary digit, and every digit you add doubles what you can name

surprise of an outcome of probability $p = -\log_2 p$



probabilities multiply, surprises add, hence the log · rare = expensive

Entropy

$$H(X) = - \sum_i p_i \log_2 p_i$$

entropy = expected surprise, each outcome weighted by how often it happens

Source coding theorem (informal): lossless codes can average $H(X)$ bits per symbol, never fewer. (Shannon 1948)

$\underbrace{\text{TTTTTTTT}}_{\text{almost always T}} \longrightarrow H \approx 0$

$\underbrace{\text{XTTXXXT}}_{\text{fair coin}} \longrightarrow 1 \text{ bit/symbol}$

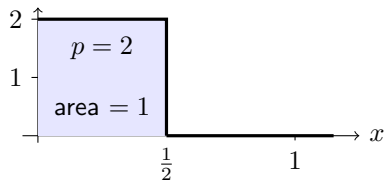
English: $4.7 \longrightarrow 4.14 \longrightarrow \approx 1 \text{ bit per letter}$

measured by people guessing text one character at a time (Shannon 1951)

Differential entropy

$$h(X) = - \int p \ln p$$

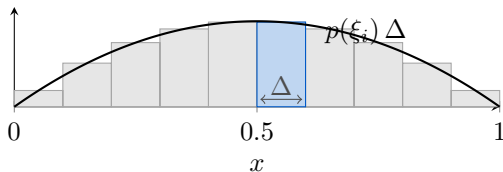
the same formula, sum to integral; areas under p are probabilities. Natural log from here on.



$$h = -\ln 2 \approx -0.69 < 0$$

$-\ln p$ is the constant $-\ln 2$, and the average of a constant is that constant
Expected surprise can't go negative. So what is h measuring?

Measuring at resolution Δ



the bins make a discrete distribution, bin i carries the shaded area $p(\xi_i)\Delta$

$H_\Delta \approx h + \ln \frac{1}{\Delta}$ the entropy of the bins (split: paper §5.1)

$\log_2 \frac{1}{\Delta} =$ binary digits to name the bin $\rightarrow \infty$ halve Δ , add one digit

$H_\Delta \rightarrow \infty$: **a continuous variable has no finite entropy.** h is the finite part.
What is it measuring?

What h is

$$h = - \int p \ln p = - \int p \ln \frac{p}{m}, \quad m \equiv 1$$

p compared, point by point, against the flat density: height one everywhere. Not an amount, a comparison.

$h > 0$: more spread than unit scale $h < 0$: more concentrated

the uniform on $[0, 1]$ sits exactly at $h = 0$. And $-\ln 2$ means: one bit cheaper than the unit interval, at every resolution.

The flat reference

$$\int m = \infty : \quad m \text{ is not a probability distribution}$$

so the comparison is not forced ≥ 0 , and the divergent $\ln \frac{1}{\Delta}$ was the reference's infinity, never p 's

h was never a continuous H . It is p measured against the flat reference.

Relative entropy

$$\sum_i p(\xi_i) \Delta \ln \frac{p(\xi_i) \Delta}{q(\xi_i) \Delta} \longrightarrow \int p \ln \frac{p}{q} = D(p \parallel q)$$

two real densities at one resolution: the Δ s cancel before any limit. No $\ln \frac{1}{\Delta}$ to remove, no correction.

$$D(p \parallel q) \geq 0, \quad = 0 \text{ only when } q = p$$

nonnegative for a genuine density q , and zero exactly when there is nothing to compare

The fundamental object

discrete was the same: $H(p) = \ln n - D(p \parallel u)$

reference minus comparison, the same split. $\ln n$ is finite, so on an alphabet it hides inside one number.

Relative entropy is the fundamental object. H is the case $q = u$; h is the case $q = m \equiv 1$.

The heat equation

$$\partial_t p = \frac{1}{2} \partial_x^2 p$$

peaks sink, valleys fill: ink spreading in water, heat evening out along a bar

$$\text{point mass} \longrightarrow \mathcal{N}(0, t), \quad h(\mathcal{N}(0, t)) = \frac{1}{2} \ln(2\pi e t)$$

the heat kernel: width $\sqrt{t}$, spreading forever, entropy growing

$$X_t = X_0 + \sqrt{t} Z$$

running the flow for time t = adding Gaussian noise of variance t . Always growing? At what rate?

de Bruijn's identity

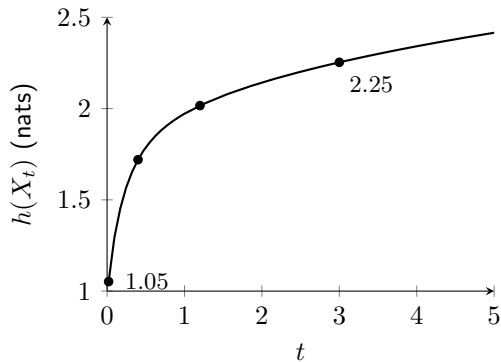
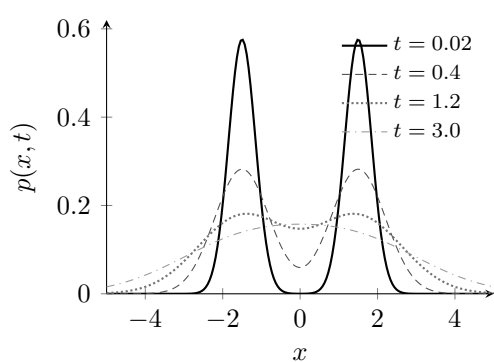
$$J(p) = \int \frac{(p')^2}{p} \quad \text{sharp features} \rightarrow \text{large } J$$

$$\boxed{\frac{d}{dt} h(p_t) = \frac{1}{2} J(p_t)}$$

Stam (1959), crediting de Bruijn

$J \geq 0$, so entropy never falls. Fast while sharp features remain, slower as the density smooths.

The picture



the slope on the right, at every instant, is $\frac{1}{2}J$ of the density on the left