

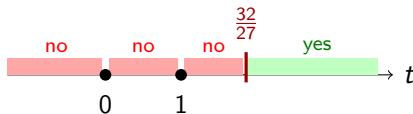
The Number $\frac{32}{27}$: A Mathematical Mystery

Chromatic Polynomials and Their Roots

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Where We Are Going



Red zones: something can **never** happen here. The sharp boundary between red and green is $\frac{32}{27} \approx 1.185$.

By the end of this talk you will know exactly what this means and why it is true.



- 1 Graphs and Colourings
- 2 The Chromatic Polynomial
- 3 Deletion-Contraction
- 4 Chromatic Roots and Tutte's Theorem
- 5 Jackson's Theorem and $32/27$
- 6 Further Results and Open Questions

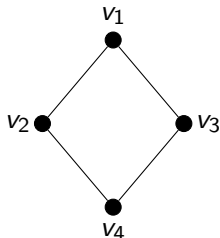


What is a Graph?

Definition

A **graph** $G = (V, E)$ has vertices V and edges E , where each edge is an unordered pair of distinct vertices.

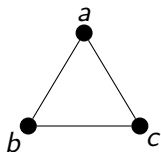
- $n = |V(G)|$ = number of vertices
- **Adjacent** = joined by an edge
- **Leaf** = vertex with exactly one edge
- **Bridge** = edge whose removal disconnects G



Three Key Graphs

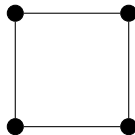


Triangle K_3



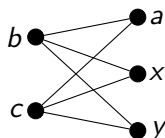
Every pair adjacent

Square C_4



Ring of 4

$K_{2,3}$



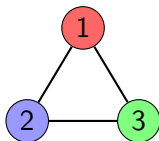
Two groups

Proper Colourings



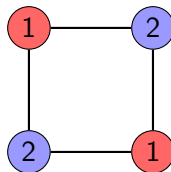
Definition

A **proper t -colouring** assigns one of t colours to every vertex so adjacent vertices never share a colour. The **chromatic number** $\chi(G)$ is the smallest such t .



$$\chi(K_3) = 3$$

2 colours fail: c touches both



$$\chi(C_4) = 2$$

Opposite vertices not adjacent

Counting Colourings: $P(G, t)$



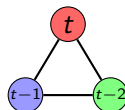
Definition

$P(G, t)$ = number of distinct proper t -colourings of G .

Compute $P(K_3, t)$:

- Colour a : t choices
- Colour b : $t - 1$ choices
- Colour c : $t - 2$ choices

$$P(K_3, t) = t(t - 1)(t - 2)$$



t	P
1	0
2	0
3	6

Key shift

Polynomial in t : plug in any real t .

It Is Always a Polynomial



Theorem

For every simple graph G on n vertices, $P(G, t)$ is a polynomial of degree n with leading coefficient 1.

Key idea. Same-colour vertices form an **independent set** (no edges between them). Let e_k = ways to split $V(G)$ into k nonempty independent sets. Then:

$$P(G, t) = \sum_{k=1}^n e_k \cdot t(t-1) \cdots (t-k+1)$$

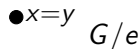
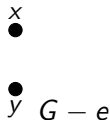
Each term is a polynomial in t . A finite sum of polynomials is a polynomial. Since $e_n = 1$, the leading term is t^n . □

Deletion-Contraction



Two operations on edge $e = xy$

$G - e$: delete the edge G/e : merge x and y



Theorem (Deletion-Contraction)

$$P(G, t) = P(G - e, t) + P(G/e, t)$$

Proof. Split colourings of $G - e$:

- x, y **differ**: valid for G . Count = $P(G, t)$.
- x, y **match**: equivalent to G/e . Count = $P(G/e, t)$.

So $P(G - e, t) = P(G, t) + P(G/e, t)$. Rearrange.

Chromatic Roots



Definitions

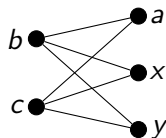
A **chromatic root** is a value of t where $P(G, t) = 0$.

$t = 0$ and $t = 1$ are **trivial** roots.

All others are **nontrivial**.

A nontrivial root exists:

$P(K_{2,3}, t) =$
 $t(t-1)(t^3-5t^2+10t-7)$ The
cubic has a real root at $t \approx 1.431$.



$K_{2,3}$: root at ≈ 1.431

Central question

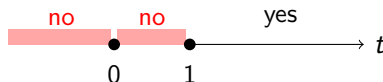
Can nontrivial roots get close to 1?

Tutte's Theorem



Theorem (Tutte, 1970s)

No graph has a chromatic root in $(-\infty, 0)$ or in $(0, 1)$.



Proof idea. Track the sign of $(-1)^{n+c}P(G, t)$ through deletion-contraction induction. This sign is permanently positive in both forbidden regions, so $P(G, t)$ can never be zero there.

Consequence: every nontrivial root satisfies $t \geq 1$.

Is There Another Gap Above 1?



Tutte: every nontrivial root satisfies $t \geq 1$.

$K_{2,3}$ has a root at ≈ 1.431 , so roots DO exist above 1.

Can roots get arbitrarily close to 1 from above?

Tutte raised this in 1974 but left it open.

Jackson's answer (1993)

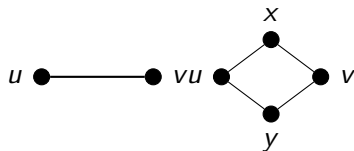
No. There is another forbidden zone above 1, with sharp boundary $\frac{32}{27} \approx 1.185$.

Double Subdivision and Generalized Triangles



Double Subdivision of uv

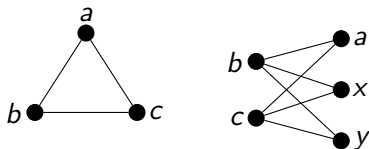
Delete uv . Add vertices x, y .
Add edges ux, xv, uy, yv .



Generalized Triangle

K_3 , or any graph obtained from K_3 by repeated double subdivisions.

One subdivision of K_3 gives $K_{2,3}$:



Roots Converging to $\frac{32}{27}$



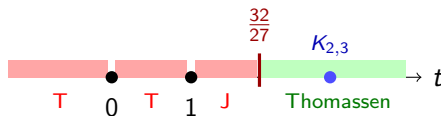
Graph	$ V $	Smallest root
K_3	3	none in $(1, 2)$
$G_1 = K_{2,3}$	5	≈ 1.4302
G_2	7	≈ 1.3611
G_3	9	≈ 1.3351
$\vdots$		$\downarrow$ $\rightarrow \frac{32}{27}$



Theorem (Jackson, 1993)

No graph has a chromatic root in $(1, \frac{32}{27})$. This bound is sharp.

The Complete Picture



- $(-\infty, 0)$ and $(0, 1)$: no roots
- $(1, \frac{32}{27})$: no roots
- $(\frac{32}{27}, \infty)$: roots dense, no further gaps

[T-Tutte]

[J-Jackson]

[Thomassen 1997]



Dong-Koh (2010):

Forbidden minor	Threshold
$K_3, K_4 - e, K_{2,3}$	2
K_4	$\frac{32}{27}$
$K_{2,4}$	≈ 1.431

Thomassen (2000):

Graphs with a Hamiltonian path have no root below $t_0 \approx 1.296$.

Open problems:

- 1 Perrett's Conjecture: threshold $> \frac{32}{27}$ iff family omits some generalized triangle.
- 2 Does $\frac{5}{4}$ threshold extend to all 3-connected graphs?
- 3 Complex roots connect to the Potts model in statistical physics.
- 4 Analogous threshold for flow polynomials?



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- [P] T. Perrett, *Chromatic roots and minor-closed families*, arXiv:1603.08255 (2018).



Scan for the animation

"Mathematics is the art of giving the same name to different things." – **Henri Poincaré**