

The Erdős–Kac Theorem

Vikram Sarkar

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The **density** of any subset $A \subseteq \mathbb{N}$ is defined to be

$$\lim_{N \rightarrow \infty} \frac{|A \cap \{1, 2, \dots, N\}|}{N},$$

if that the limit exists. (Otherwise A does not have a density.)

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We say $g(n) \ll f(n)$ if $g(n) = O(f(n))$. For two fixed constants a and b , we say $a \ll b$ instead of $a < b$ to indicate that a is “sufficiently less than b .”

Definition of $\omega(n)$

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Question

How big is $\omega(n)$? That is, roughly speaking, how large do we expect it to be given, say, n is some fixed order of magnitude? (More rigorously speaking, what is an “average order” of $\omega(n)$?)

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Definition

This won't be very central to the main topic, but just to make this notion of “largeness” concrete: f and g have the same **average order** if

$$\lim_{x \rightarrow \infty} \frac{\sum_{n \leq x} f(n)}{\sum_{n \leq x} g(n)} = 1.$$

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Answer to the question

The answer is that $\omega(n) \approx \log \log n$. Meaning, they have the same average order.

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- For each prime p , let $\chi_p = \begin{cases} 1 & p \mid n \\ 0 & p \nmid n \end{cases}$.
- Then, $\omega(n) = \sum_{p \leq n} \chi_p$ since $\chi_p = 1$ iff $p \mid n$.

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- So, we can “approximate” χ_p as $\frac{1}{p}$. Then, $\omega(n) \approx \sum_{p \leq n} \frac{1}{p}$.

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Theorem (Mertens, 1874)

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We can use this logic to prove that $\log \log n$ and $\omega(n)$ have the same average order. (Exercise for the reader.)

Further Estimates of $\omega(n)$

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Theorem (Hardy and Ramanujan, 1917)

The **normal order** of $\omega(n)$ is $\log \log n$, that is, for all $\varepsilon > 0$, that is, the density of all $n \in \mathbb{N}$ for which

$$\left| \frac{\omega(n)}{\log \log n} - 1 \right| < \varepsilon$$

is 1.

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Theorem (Turán, 1934)

$$\sum_{n \leq x} (\omega(n) - \log \log n)^2 = (1 + o(1))x \log \log x.$$

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- In fact, the error term has a strong connection with $\sqrt{\log \log n}$:

Theorem (Erdős, Kac, 1959)

If $a < b$ are in $\mathbb{R} \cup \{-\infty, \infty\}$, the density of the set

$$\left\{ n \in \mathbb{N} \mid a \leq \frac{\omega(n) - \log \log n}{\sqrt{\log \log n}} \leq b \right\}$$

is

$$\frac{1}{\sqrt{2\pi}} \int_a^b e^{-x^2/2} \, dx.$$

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Okay, what?

Small Simplification

It is not hard to show that it suffices to show the density of the set

$$\left\{ n \in \mathbb{N} \mid a \leq \frac{\omega(n) - \log \log N}{\sqrt{\log \log N}} \leq b \right\}$$

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instead. Since $\log \log N = \sum_{p \leq N} p^{-1} + O(1)$, it is also not hard to show that if $A(N) = \sum_{p \leq N} p^{-1}$, it suffices to show that the density of the set

$$\left\{ n \in \mathbb{N} \mid a \leq \frac{\omega(n) - A(N)}{A(N)^{1/2}} \leq b \right\}.$$

is $\frac{1}{\sqrt{2\pi}} \int_a^b e^{-x^2/2} dx$ as well.

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- A graph of φ is shown below.

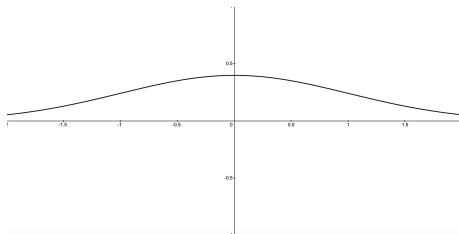
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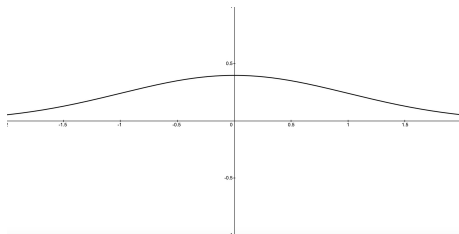
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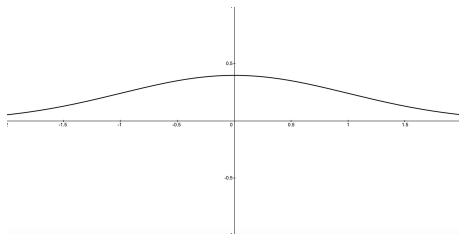
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- Looks like a “squished” bell curve.
- Note that $\int_{-\infty}^{\infty} \varphi(x) \, dx = 1$ due to the infamous $\int_{-\infty}^{\infty} e^{-x^2} \, dx = \sqrt{\pi}$.

Probability Theory (Continued)

Definition

A **discrete random variable** X has a finite sample space S and fixed probabilities of being each element of S . For example, S could be $\{1, 2, 3\}$, and $\text{Prob}(X = 1) = \frac{1}{3}$, $\text{Prob}(X = 2) = \frac{1}{6}$, $\text{Prob}(X = 3) = \frac{1}{2}$.

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Definition

A **continuous random variable** X which has an infinite sample space $S \subseteq \mathbb{R}$ that is equipped with a **probability density function** f has

$$\text{Prob}(a \leq X \leq b) = \int_a^b f(x)dx.$$

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In our situation, we have the sequence of discrete random variables $\{X_N\}$, for which X_N has sample space $\left\{ \frac{\omega(n) - A(N)}{A(N)^{1/2}} \mid n \leq N \right\}$, and each $n \leq N$ has equal probability to be chosen.

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$$\mathbb{E}[X] := \frac{1}{|S|} \sum_{s \in S} s \cdot \text{Prob}(X = s).$$

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$$\mathbb{E}[X^k] := \frac{1}{|S|^k} \sum_{s \in S} s^k \cdot \text{Prob}(X = s).$$

Moments (Continued)

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Theorem (Well-known)

The standard normal distribution is completely determined by all its moments. That is, if there exists some random variable X with

$$\mathbb{E}[X^k] = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} x^k e^{-x^2/2} dx,$$

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Theorem (Well-known)

Suppose $X_1, X_2, \dots, X_n, \dots$ is a sequence of random variables. Let X be a random variable that is completely determined by all its moments. Then, if $\mathbb{E}[X_n^k] \rightarrow \mathbb{E}[X^k]$ for all $k \geq 1$, then $\text{Prob}(a \leq X_n \leq b) \rightarrow \text{Prob}(a \leq X \leq b)$ for all $a < b \in \mathbb{R} \cup \{-\infty, \infty\}$.

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- Now, let X_N be the random variable with sample space $\left\{ \frac{\omega(n) - A(N)}{A(N)^{1/2}} \mid n \leq N \right\}$ for which each element of this set has an equal probability of being chosen (if there are somehow repeats, then the probability of $X_n = r$ would be the # of n for which $\frac{\omega(n) - A(N)}{\sqrt{A(N)}} = r$ divided by N).

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- Let X be a random variable with probability density function φ . Then, if we had

$$\mathbb{E}[X_N^k] \rightarrow \mathbb{E}[X^k] = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} x^k e^{-x^2/2} dx,$$

we would then have

$\text{Prob}(a \leq X_N \leq b) \rightarrow \text{Prob}(a \leq X \leq b) = \frac{1}{\sqrt{2\pi}} \int_a^b e^{-x^2/2} dx$. where X is a random variable with probability density function φ .

Moments (Continued)

- Now, let X_N be the random variable with sample space $\left\{ \frac{\omega(n)-A(N)}{\sqrt{A(N)}} \mid n \leq N \right\}$ for which each element of this set has an equal probability of being chosen (if there are somehow repeats, then the probability of $X_n = r$ would be the $\#$ of n for which $\frac{\omega(n)-A(N)}{\sqrt{A(N)}} = r$ divided by N).
- Let X be a random variable with probability density function φ . Then, if we had

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$\text{Prob}(a \leq X_N \leq b) \rightarrow \text{Prob}(a \leq X \leq b) = \frac{1}{\sqrt{2\pi}} \int_a^b e^{-x^2/2} dx$. where X is a random variable with probability density function φ .

- However, note that $\lim_{N \rightarrow \infty} (a \leq X_N \leq b)$ is also the density of $\{n \in \mathbb{N} \mid \frac{\omega(n)-A(N)}{\sqrt{A(N)}} \in [a, b]\}$. So showing that the moments of X_N approach the moments of X suffices.

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Theorem

We have

$$\mu_k := \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\infty} x^k e^{-x^2/2} = \begin{cases} (k-1)!! & k \text{ even}, k > 2 \\ 0 & k \text{ odd} \\ 1 & k = 0 \end{cases}.$$

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This is not very hard to prove; if k is odd then we are integrating a bounded odd function over a symmetric interval so the integral is 0. If k is even, use integration by parts to get a recurrence between μ_k and μ_{k+2} . So, we want to prove $\mathbb{E}[X_N^k] \rightarrow \mu_k$.

Simplifications

- Note that

$$\mathbb{E}[X_N^k] = \frac{1}{N} \sum_{n \leq N} \left(\frac{\omega(n) - A(N)}{A(N)^{1/2}} \right)^k.$$

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- So, we want to show that

$$\frac{1}{N} \sum_{n \leq N} \left(\frac{\omega(n) - A(N)}{A(N)^{1/2}} \right)^k = \mu_k + o(1).$$

$$L_p(n)$$

Upon rearrangement, we want to show that

$$\sum_{n \leq N} (\omega(n) - A(N))^k = N(A(N))^{k/2} \mu_k(1 + o(1)).$$

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The important thing here is that $\sum_{p \leq N} L_p(n) = \omega(n) - A(N)$.

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Therefore, we wish to show that

$$\sum_{n \leq N} \left(\sum_{p \leq N} L_p(n) \right)^k = N(A(N))^{k/2} \mu_k(1 + o(1)).$$

$$L(p)$$

Now, for each prime p , let $L(p)$ be a random variable for which $L(p) = -\frac{1}{p}$ with probability $\frac{1}{p}$ and $L(p) = 1 - \frac{1}{p}$ with probability $1 - \frac{1}{p}$.

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Now, for each prime p , let $L(p)$ be a random variable for which $L(p) = -\frac{1}{p}$ with probability $\frac{1}{p}$ and $L(p) = 1 - \frac{1}{p}$ with probability $1 - \frac{1}{p}$. The following two theorems are not very hard to prove, for the sake of time I won't be going over their proofs.

Theorem

Suppose $N \geq 1$, and $N_S = N^{1/\log(\sqrt{\log \log N} + 3)}$. Then,

$$\sum_{p \leq N} L_p(n) = \sum_{p \leq N_S} L_p(n) + O(\log(\sqrt{\log \log N} + 3)),$$

for all $n \leq N$.

Relating L_p to $L(p)$

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Theorem

Let $\pi(n)$ be the number of primes $\leq n$. Then, for any positive integer k , we have

$$\sum_{n \leq N} \left(\sum_{p \leq N_S} L_p(n) \right)^k = N \mathbb{E} \left[\left(\sum_{p \leq N_S} L(p) \right)^k \right] + O(3^k \pi(N_S)^k)$$

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So we have expressed something very similar to what we want (the LHS) as something relating the random variable $L(p)$. Right now, the right hand side is not very useful since the main expression is stuck in an expected value. So we need a good way to approximate it.

Moment Generating Functions

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Proof

It is clear that M_X is analytic on \mathbb{C} . Note that from the Taylor Series expansion of e^x ,

$$M_X(z) = \mathbb{E}[e^{zX}] = \mathbb{E} \left[\sum_{k \geq 0} \frac{(zX)^k}{k!} \right] = \sum_{k \geq 0} \mathbb{E} \left[\frac{(zX)^k}{k!} \right] = \sum_{k \geq 0} \frac{\mathbb{E}[X^k]}{k!} z^k.$$

Thus $M_X^{(k)}(0) = \mathbb{E}[X^k]$, as desired.

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$$\begin{aligned} M_{B_N}(t) &= \mathbb{E} \left[e^{t \sum_{p \leq N} L(p)} \right] \\ &= \prod_{p \leq N} \mathbb{E} \left[e^{tL(p)} \right]. \end{aligned}$$

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- Now, for $|t| \leq \frac{1}{2}$, we can say by Taylor Series that

$$e^{tL(p)} = 1 + tL(p) + \frac{t^2}{2} L(p)^2 + O(|t|^3 |L(p)|^3),$$

so

$$\mathbb{E}[e^{tL(p)}] = 1 + \frac{p-1}{2p^2} t^2 + O\left(\frac{|t|^3}{p}\right).$$

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- One can then take the product over all $p \leq N$ and then bound to obtain that

$$M_{B_N} \left(\frac{t}{\sqrt{\log \log N}} \right) \rightarrow e^{t^2/2}$$

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- Then we can compute

$$\mathbb{E}[B_N^k] = M_{B_N}^{(k)}(0) = (1 + o(1))\mu_k(\log \log N)^{k/2},$$

since the k^{th} derivative of $e^{x^2/2}$ at 0 is μ_k (!!) from the Taylor Series expansion of e^x .

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Theorem

$$\sum_{n \leq N} \left(\sum_{p \leq N} L_p(n) \right)^k - \sum_{n \leq N} \left(\sum_{p \leq N_S} L_p(n) \right)^k = o(N(\log \log N)^{k/2}).$$

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Now, we can use the result we obtained above to get the sum for $p \leq z$.

The Finish (Continued)

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Note that

$$\begin{aligned}(\log \log N)^{1/2} - (\log \log N_S)^{1/2} &\leq (\log \log N - \log \log N_S)^{1/2} \\&= \left(\log \log N - \log \left(\frac{\log N}{\log(\sqrt{\log \log N} + 3)} \right) \right)^{1/2} \\&= (\log \log(\sqrt{\log \log N} + 3))^{1/2}\end{aligned}$$

The Finish (Continued)

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so by the Binomial Theorem,

$$(\log \log N)^{k/2} - (\log \log N_S)^{k/2} = O \left(k (\log \log N_S)^{\frac{k-1}{2}} (\log \log(\sqrt{\log \log N} + 3)) \right),$$

The Finale

Putting it altogether, we have

$$\begin{aligned}\sum_{n \leq N} \left(\sum_{p \leq N} L_p(n) \right)^k &= (1+o(1)) N (\log \log z)^{k/2} \mu_k + O(3^k \pi(N_S)^k) \\ &\quad + o(N (\log \log N)^{k/2}) \\ &= (1+o(1)) N (\log \log N)^{k/2} \mu_k + O(k (\log \log z)^{\frac{k-1}{2}}) + \dots \\ &= (1+o(1)) N (\log \log N)^{k/2} \mu_k \\ &= (1+o(1)) N (A(N))^{k/2} \mu_k,\end{aligned}$$

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so we are done. ■

Thanks!

Any questions?