Shor's Algorithm: A Quantum Leap in Factorization Lets Decrypt the Algorithm

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- 2. Context and Importance
- 3. Background: Quantum Computing Basics
- 4. Core Concept: Period Finding
- 5. Key Steps of Shor's Algorithm
- 6. Conclusion

The Cryptographic Landscape

- As of 2021:
 - 52% of HTTPS servers use RSA
 - 75% of digital certificates use RSA
- Safeguarding trillions of online transactions
- But a quantum storm is brewing...



Image Source: https://threatpost.com/why-web-browserpadlocks-shouldnt-be-trusted/159659/

Key Point

RSA encryption is the backbone of current internet security.

The Quantum Revolution

- Quantum computing is evolving rapidly
- Quantum volumes increasing 10 fold yearly since 2020
- Potential to completely shift the cryptographic landscape



Image: Exponential growth of quantum volume (https://www.quantinuum.com/news/quantinuumextends-its-significant-lead-in-quantum-computingachieving-historic-milestones-for-hardware-fidelity-andquantum-volume)

The Problem of Integer Factorization

- Central problem in number theory and computer science
- Difficulty increases exponentially with number size
- Example: Factoring a 2048-bit number
 - Classical computers: Billions of years
 - Quantum computers with Shor's algorithm: Hours or days
- Forms the foundation of many cryptographic systems, especially RSA

Definition: Integer Factorization

The process of decomposing a composite number into a product of smaller integers.

Peter Shor and His Algorithm

- Peter Shor: American mathematician and MIT professor
- Developed Shor's algorithm in 1994 at AT&T Bell Laboratories
- One of the algorithms to show quantum computers could exponentially outperform classical computers on a problem of wide interest
- Sparked intense interest in quantum computing and quantum-resistant cryptography



Image: Peter Shor

The Breakthrough: Polynomial-Time Factorization

- Shor's algorithm: Integer factorization in polynomial time on a quantum computer
- Dramatic improvement over classical methods
- Can factor an n-bit number in $O(n^3)$ time and O(n) space
- Implications:
 - Many current cryptographic systems will be vulnerable to attack
 - Need for quantum-resistant cryptography

Key Insight

Shor's algorithm reduces factoring to finding the period of a quantum function.

Qubits: The Fundamental Unit

- Qubit: Quantum bit, the basic unit of quantum information
- Unlike classical bits, qubits can be in superposition
- Represented as $|\psi\rangle = \alpha |\mathbf{0}\rangle + \beta |\mathbf{1}\rangle$
- α and β are complex numbers: $|\alpha|^2 + |\beta|^2 = 1$



Bloch sphere representation of a qubit

Superposition Principle

A qubit can exist in a superposition of multiple states until measured.

Quantum State Representation

- Quantum states are represented by vectors in complex Hilbert space
- For a single qubit:

$$\ket{0} = egin{pmatrix} 1 \ 0 \end{pmatrix}, \quad \ket{1} = egin{pmatrix} 0 \ 1 \end{pmatrix}$$

• General state:

$$|\psi
angle = lpha |\mathbf{0}
angle + eta |\mathbf{1}
angle = egin{pmatrix} lpha \ eta \end{pmatrix}$$

• Multiple qubits: tensor product of individual qubit states

Key Point

The state space grows exponentially with the number of qubits!

Tensor Product: Combining Quantum Systems

- Tensor product () combines individual qubit states
- For two qubits $|\psi_1\rangle = a|0\rangle + b|1\rangle$ and $|\psi_2\rangle = c|0\rangle + d|1\rangle$:

 $|\psi_1
angle\otimes|\psi_2
angle=ac|00
angle+ad|01
angle+bc|10
angle+bd|11
angle$

• State space grows exponentially: n qubits require 2ⁿ amplitudes

Key Point

Tensor product enables description of multi-qubit systems!



Quantum Fourier Transform (QFT)

- Quantum analogue of the classical Fourier transform
- Crucial component in many quantum algorithms, including Shor's
- Transforms quantum state from computational basis to Fourier basis
- For an n-qubit state $|x\rangle$:

$$QFT|x
angle=rac{1}{\sqrt{2^n}}\sum_{y=0}^{2^n-1}e^{2\pi ixy/2^n}|y
angle$$

• Can be implemented efficiently using $O(n^2)$ quantum gates

QFT Superpower

QFT can extract periodicity information from quantum states!

Quantum Parallelism and Interference

Quantum Parallelism:

- Ability to perform operations on many computational states at the same time
- Enabled by superposition
- Example: Evaluating a function for multiple inputs

Quantum Interference:

- Amplitudes can interfere constructively or destructively
- Crucial for extracting useful information from quantum computations



Quantum interference

Key Insight

Quantum parallelism and interference are key to quantum speedups!

Reducing Factoring to Period Finding

- Key insight: Factoring can be reduced to finding the period of a function
- For a number *N* to be factored, define:

$$f(x) = a^x \mod N$$

where a is coprime to N

- This function is periodic: f(x) = f(x + r) for some r
- Finding this period r can lead to factors of N

Key Point

Period finding is hard classically but efficient quantumly!

- If we find the period *r*:
 - Compute $a^{r/2} \mod N$
 - If this equals $\pm 1 \mod N$, try next a
 - Otherwise, $gcd(a^{r/2} \pm 1, N)$ likely gives a factor
- Example: For N = 15, a = 7
 - Period r = 4
 - $7^2 \mod 15 = 4$
 - gcd(4-1,15) = 3 and gcd(4+1,15) = 5

- 1. Quantum state preparation
- 2. Modular exponentiation
- 3. Quantum Fourier Transform
- 4. Measurement and classical post-processing

Step 1: Quantum State Preparation

• Initialize two quantum registers:

- Input register: $n = 2 \lceil \log_2 N \rceil$ qubits
- Output register: $\lceil \log_2 N \rceil$ qubits
- Apply Hadamard gates to create superposition:

$$|\psi_1
angle = rac{1}{\sqrt{2^n}}\sum_{x=0}^{2^n-1}|x
angle|0
angle$$

Key Point

This superposition allows us to evaluate the function for all inputs simultaneously!

• Apply the function $f(x) = a^x \mod N$ to the superposition:

$$|\psi_2
angle = rac{1}{\sqrt{2^n}}\sum_{x=0}^{2^n-1}|x
angle|a^x \mod N
angle$$

- Implemented using controlled modular multiplication
- Most computationally intensive part of the algorithm

Step 3: Quantum Fourier Transform

• Apply QFT to the input register:

$$|\psi_3
angle = rac{1}{2^n}\sum_{y=0}^{2^n-1}\sum_{x=0}^{2^n-1}e^{2\pi i x y/2^n}|y
angle|a^x \mod N
angle$$

- Transforms periodicity in function values to phase differences
- Efficient implementation using $O(n^2)$ gates

Key Insight

QFT allows us to extract period information efficiently!

- 1. Measure the input register to obtain y
- 2. Use continued fraction expansion to find r' approximating $\frac{2^n}{r}$
- 3. Check if $a^{r'} \equiv 1 \pmod{N}$
- 4. If r' is even, compute $\gcd(a^{r'/2} \pm 1, N)$
- 5. If we find a non-trivial factor, we're done; otherwise, repeat

Success Probability

The algorithm succeeds with probability $\Omega(1/\log \log N)$ per iteration

Implications for Cryptography

- Shor's algorithm threatens RSA and other public-key cryptosystems
- Need for quantum-resistant cryptography:
 - Lattice-based cryptography
 - Hash-based signatures
 - Code-based cryptography
 - Multivariate cryptography
- NIST Post-Quantum Cryptography Standardization

Key Point

We need to prepare for a post-quantum cryptographic landscape!

Current State and Future Prospects

- Largest number factored using Shor's: 21 (as of 2012)
- Challenges:
 - Quantum error correction
 - Maintaining coherence
 - Scaling up number of qubits
- Ongoing research to improve implementation
- Potential impact beyond cryptography

Recap and Final Thoughts

- Shor's algorithm: A quantum solution to integer factorization
- Exponential speedup over classical algorithms
- Key components:
 - Quantum parallelism
 - Period finding
 - Quantum Fourier Transform
- Significant implications for cryptography and beyond
- Drives development in quantum computing and post-quantum cryptography

Final Thought

Shor's algorithm exemplifies the transformative potential of quantum computing!